Randomized Algorithms for Canonical Polyadic Decomposition

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Outline

Canonical polyadic decomposition (CPD)

Randomized algorithms for CPD

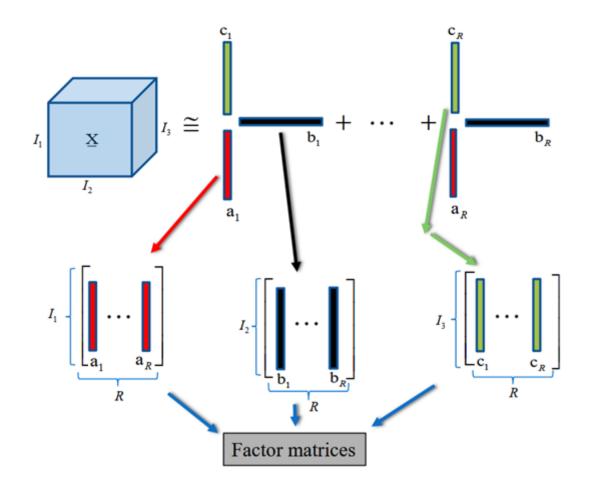
Randomized CPD with a prior reduction in the Tucker format

Randomized block sampling algorithm

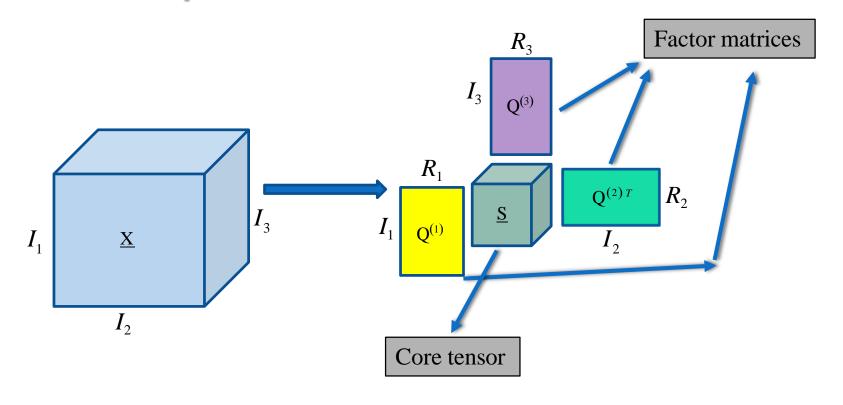
Randomized ALS algorithm

- Randomized Parafac2
- Simulations

Canonical Polyadic Decomposition (CPD)



Tucker decomposition



Orthogonal factor matrices and all-orthogonality of the core tensor HOSVD

Algorithms for CPD

- Alternating Least Squares (ALS) algorithm
- Gauss-Newton and Levenberg-Marquardt algorithms
- Tensor power iteration algorithm
- Fast damped Gauss-Newton (dGN)
- Simultaneous matrix diagonalization

Alternating least-squares algorithm (ALS)

Fixing all factor matrices and updating one by solving the following least squares problem

$$\mathbf{Min}_{\mathbf{A}^{(n)}} \| Z^{(n)} \mathbf{A}^{(n)T} - \mathbf{X}_{(n)}^{T} \|_{F}, \ n = 1, 2, ..., N$$

$$Z^{(n)} = \mathbf{A}^{(N)} \odot \cdots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \cdots \odot \mathbf{A}^{(1)}$$

Explicit solution
$$A^{(n)} = X_{(n)}Z^{(n)} W^{(n)\dagger}$$

$$\mathbf{W}^{(n)} = \mathbf{A}^{(1)T} \mathbf{A}^{(1)} * \dots * \mathbf{A}^{(n-1)T} \mathbf{A}^{(n-1)} * \mathbf{A}^{(n+1)T} \mathbf{A}^{(n+1)} * \dots * \mathbf{A}^{(N)T} \mathbf{A}^{(N)}$$

CPD with prior HOSVD compression

Computational complexity

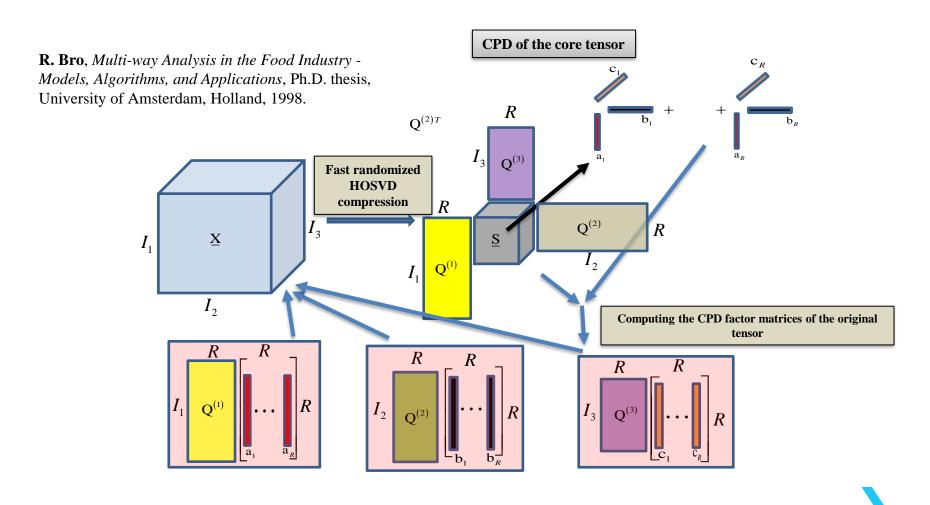
3-th order tensors
$$O\left(3R\prod_{i=1}^{3}I_{i}\right)$$
 $O\left(3R^{4}\right)$

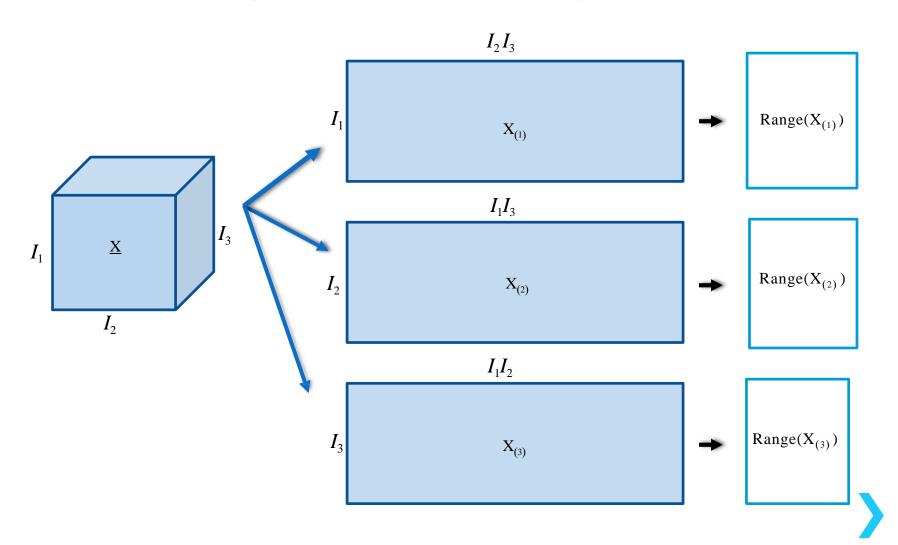
N-th order tensors
$$O\left(NR\prod_{i=1}^{N}I_{i}\right)$$
 $O\left(NR^{N+1}\right)$

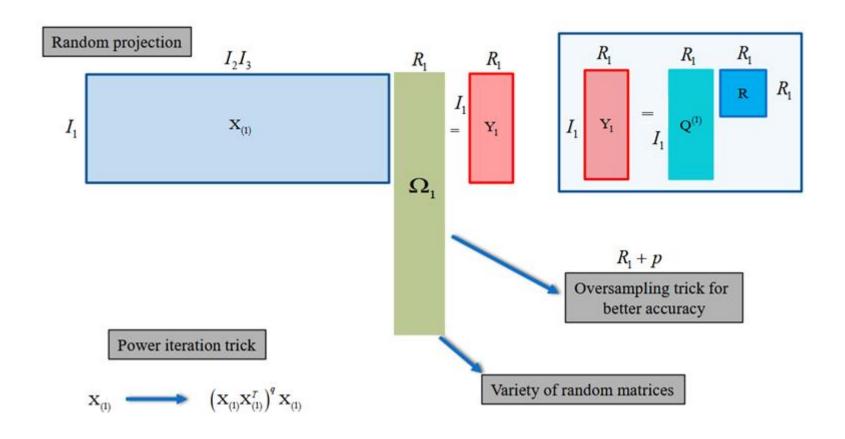
This approach is applicable when a tensor is not of very high order and also the tensor rank is less than the original tensor sizes.

Compression step can be performed by the randomized HOSVD algorithms

CPD with prior HOSVD compression







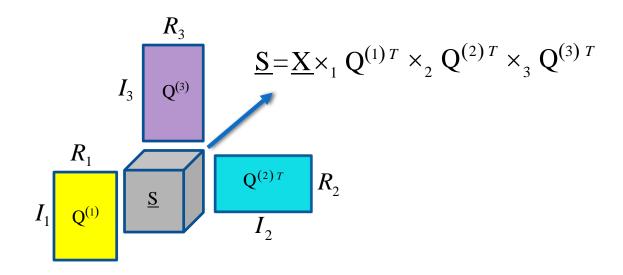
Sparse random matrices

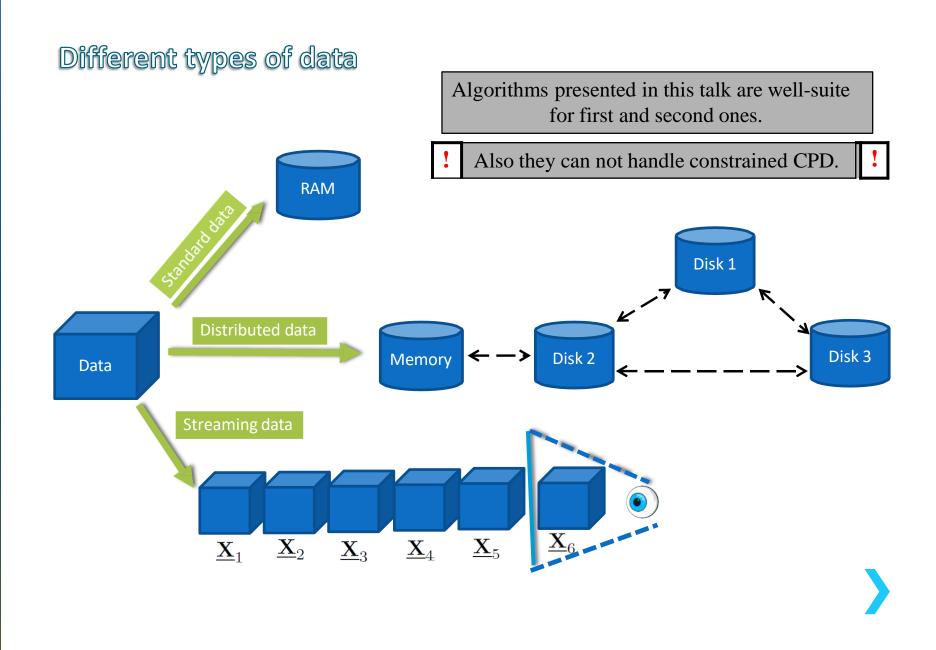
$$\Phi = \begin{cases} -1 \text{ with probability } \frac{1}{6} \\ 0 \text{ with probability } \frac{2}{3} \\ +1 \text{ with probability } \frac{1}{6} \end{cases}$$

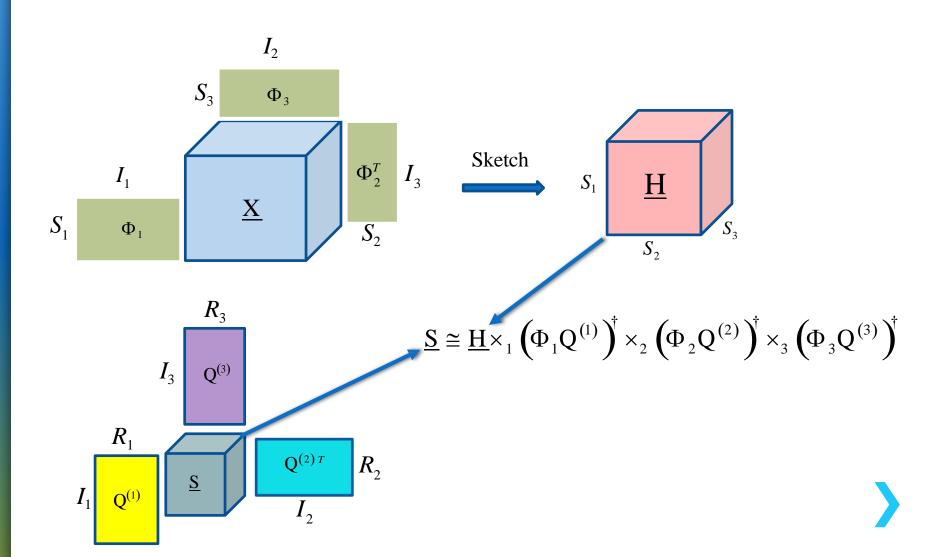
$$\Phi = \begin{cases} -1 \text{ with probability } \frac{1}{2\sqrt{D}} \\ 0 \text{ with probability } 1 - \frac{1}{\sqrt{D}} \\ +1 \text{ with probability } \frac{1}{2\sqrt{D}} \end{cases}$$

where matrix components are i.i.d.

These random matrices are reminiscent the quantization procedure in learning DNNs.



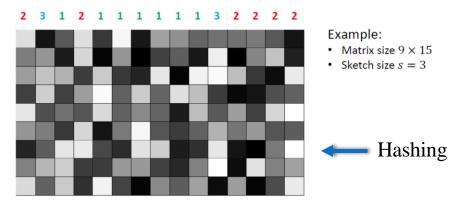




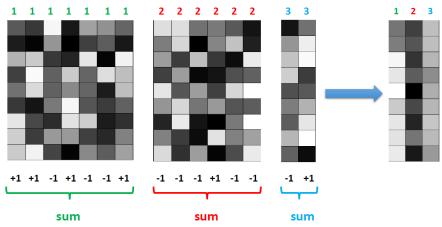
Count sketch

Pictures from:

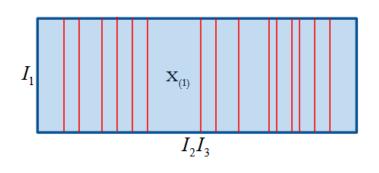
S. Wang. A practical guide to randomized matrix computations with MATLAB implementations. arXiv:1505.07570, 2015.

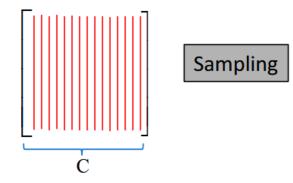


(a) Hash each column with a value uniformly sampled from $[s] = \{1, 2, 3\}$.



(b) Flip the sign of each column with probability 50%, and then sum up columns with the same hash value.





Sampling columns

Length-squared distribution

Leverage scores distribution

$$\mathbf{X}_{(n)} \in \mathbf{R}^{I_n \times \prod_{m \neq n} I_m} \qquad \qquad \begin{matrix} \mathbf{j}\text{-th column leverage score} \\ \\ l_j = \left\| \mathbf{V}_R \left(j, : \right) \right\|_2^2, \quad j = 1, 2, \dots, \prod_{m \neq n} I_m \\ \\ \\ \sum_{j=1}^{m \neq n} l_j = R \end{matrix}$$

i-th row leverage score
$$\leftarrow l_i = \left\| \mathbf{U}_R \left(i, : \right) \right\|_2^2, \quad i = 1, 2, \dots, I_n$$

$$\sum_{i=1}^{I_n} l_i = R$$

Leverage score distribution for column selection

$$p_{j} = \frac{l_{j}}{R}, \quad j = 1, 2, ..., \prod_{m \neq n} I_{m}$$

Maximum of the leverage scores is called coherence of a matrix

Fast algorithms for the computations of the leverage scores are proposed in the following papers.

Journal of Machine Learning Research 13 (2012) 3441-3472

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Fast Approximation of Matrix Coherence and Statistical Leverage

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On Fast Leverage Score Sampling and Optimal Learning

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Abstract

Leverage score sampling provides an appealing way to perform approximate computations for large matrices. Indeed, it allows to derive faithful approximations with a complexity adapted to the problem at hand. Yet, performing leverage scores sampling is a challenge in its own right requiring further approximations. In this paper, we study the problem of leverage score sampling for positive definite matrices defined by a kernel. Our contribution is twofold. First we provide a novel algorithm for leverage score sampling and second, we exploit the proposed method in statistical learning by deriving a novel solver for kernel ridge regression. Our main technical contribution is showing that the proposed algorithms are currently the most efficient and accurate for these problems.

Relative error

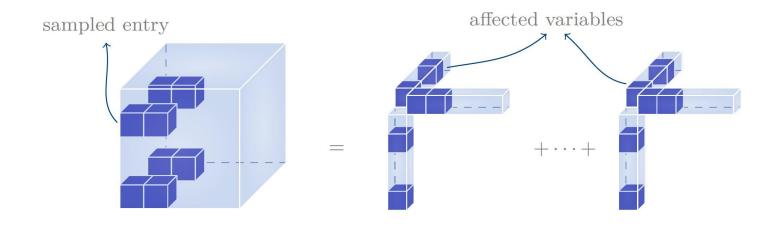
$$\|\mathbf{X} - \mathbf{C}\mathbf{C}^{\dagger}\mathbf{X}\| \leq \varepsilon \|\mathbf{X} - \mathbf{X}_{R}\|$$

$$\|\mathbf{X} - \mathbf{C}\mathbf{C}^{\dagger}\mathbf{X}\| \le \|\mathbf{X} - \mathbf{X}_{R}\| + \varepsilon \|\mathbf{X}_{R}\|$$
 Additive error

Randomized rank revealing or equivalently randomized fixed-precision algorithms are applicable when an estimation of matrix rank is unknown and it is approximated adaptively by the underlying algorithms.

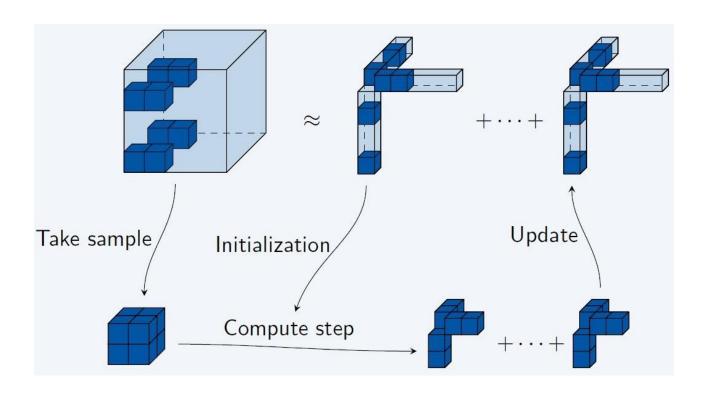
Randomized block sampling algorithm

Randomized block sampling uses the locality property to compute the CPD of a given tensor.



Picture from: N. Vervliet, *Compressed sensing approaches to large-scale tensor decompositions*, PhD thesis, KU Leuven, May 2018.

Randomized block sampling algorithm

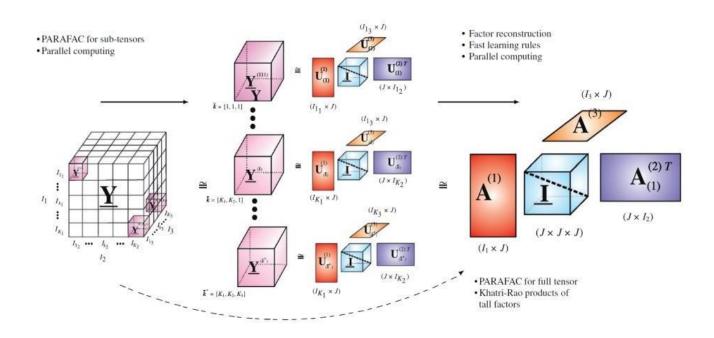


Picture from: N. Vervliet, *Compressed sensing approaches to large-scale tensor decompositions*, PhD thesis, KU Leuven, May 2018.

N. Vervliet, O. Debals, L. Sorber, M. V. Barel, L. D. Lathauwer, Tensorlab 3.0.

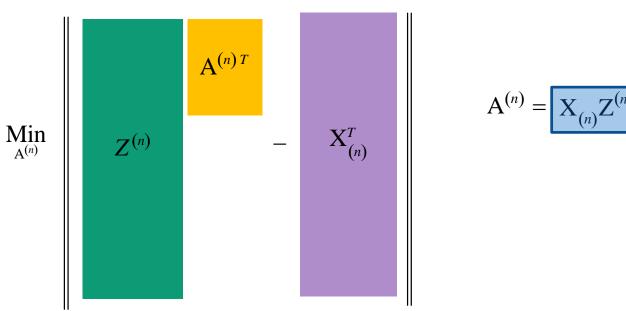
Randomized block sampling algorithm

A sampling CPD is also proposing in the following paper

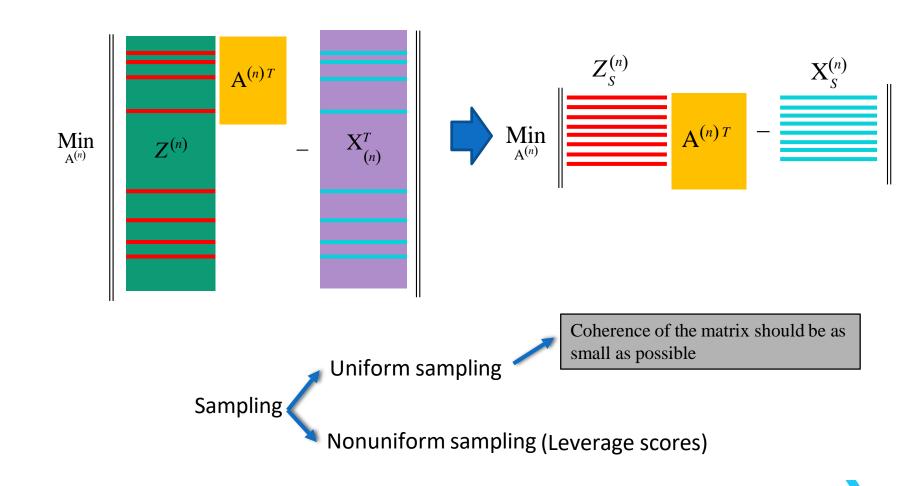


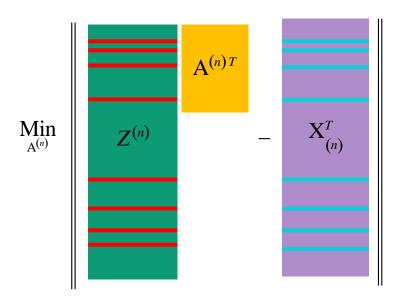
$$\mathbf{Min}_{A} \| Z^{(n)} \mathbf{A}^{(n)T} - \mathbf{X}_{(n)}^{T} \|_{F}, n = 1, 2, ..., N$$

$$Z^{(n)} = \mathbf{A}^{(N)} \odot \cdots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \cdots \odot \mathbf{A}^{(1)}$$



$$\mathbf{A}^{(n)} = \mathbf{X}_{(n)} \mathbf{Z}^{(n)} \mathbf{W}^{(n) \dagger}$$





The row selection does not need computing $Z^{(n)}$ explicitly.

$$Z^{(n)}(j,:) = A^{(1)}(i_1,:) * \cdots * A^{(n-1)}(i_{n-1},:) * A^{(n+1)}(i_{n+1},:) * \cdots * A^{(N)}(i_N,:)$$

$$j = 1 + \sum_{k=1, k \neq n}^{N} (i_k - 1) J_k, \quad J_k = \prod_{m=1, m \neq n}^{N-1} I_m$$

The nonuniform sampling is used in the following paper.

SPALS: Fast Alternating Least Squares via Implicit Leverage Scores Sampling



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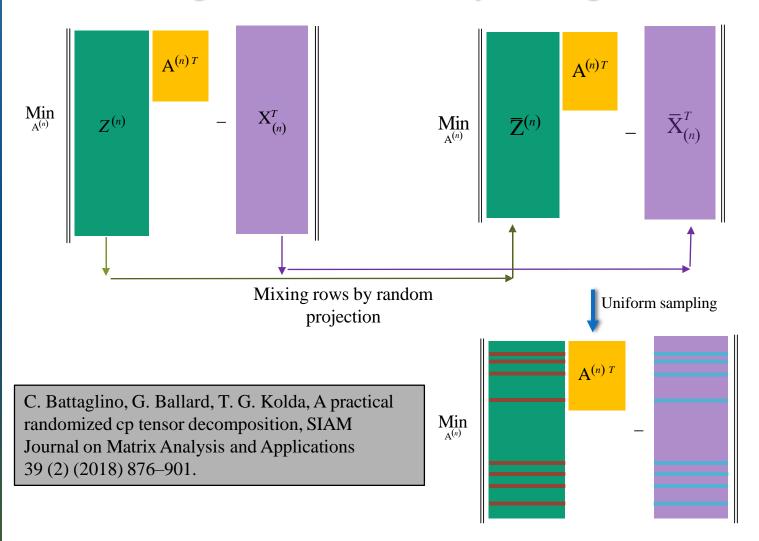
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For both Khatri-Rao and Kroncker product of matrices



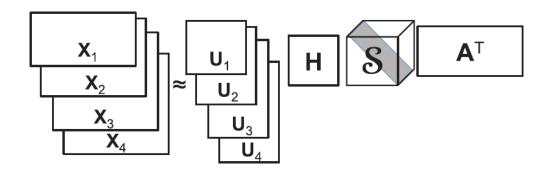


Matlab codes are included in the tensor toolbox (Matlab) and Tensorly (Python)

- **B. W. Bader, T. G. Kolda**, Matlab tensor toolbox version 2.6 (2015).
- **J. Kossaifi, Y. Panagakis, A. Anandkumar, M. Pantic, Tensorly**: *Tensor learning in python*, The Journal of Machine Learning Research 20 (1) (2019) 925–930.

Randomized algorithms for parafac2

Parafac2 model



Picture from: E.E. Papalexakis, C. Faloutsos, N.D. Sidiropoulos, *Tensors for data mining and data fusion: Models, applications, and scalable algorithms*, ACM Trans. on Intelligent Systems and Technology, vol. 8, no. 2, 16:1–16:44, 2016.

Randomized algorithms for parafac2

```
ALGORITHM 8: ALS Algorithm for PARAFAC2
Input: Multiset \{X_k\} for k = 1 : K and rank R.
Output: PARAFAC2 Decomposition of \{X_k\}: \{U_k\}, H, S, V.
 1: Initialize:
     \mathbf{V} \leftarrow R \text{ principal eigenvectors of } \sum \mathbf{X}_k^T \mathbf{X}_k
    \mathbf{H} \leftarrow \mathbf{I}
                                                                                                          A variety of randomized
 2: for k = 1 \cdots K do
      S(:,:,k) \leftarrow \mathbf{I}, for k = 1 \cdots K.
                                                                                                          algorithms can be used.
 4: end for
 5: while convergence criterion is not met do
        for k = 1 \cdots K do
           [\mathbf{P}_k, \mathbf{\Sigma}_k, \mathbf{Q}_k] \leftarrow \text{truncated SVD of } \mathbf{H} \mathbf{S}_k \mathbf{V}^T \mathbf{X}_k^T \text{ at rank } R
           \mathbf{U}_k \leftarrow \mathbf{Q}_k \mathbf{P}_k^T
 8:
        end for
 9:
        for k = 1 \cdots K do
10:
            Compute \mathcal{Y}(:,:,k) = \mathbf{U}_{b}^{T} \mathbf{X}_{k}
11:
         end for
12:
        Run a single iteration of CP ALS (Algorithm 1) on y and compute factors H, V, Ŝ.
13:
        for k = 1 \cdots K do
14:
            S(:,:,k) \leftarrow Diag(\hat{\mathbf{S}}(k,:))
15:
16:
         end for
17: end while
```

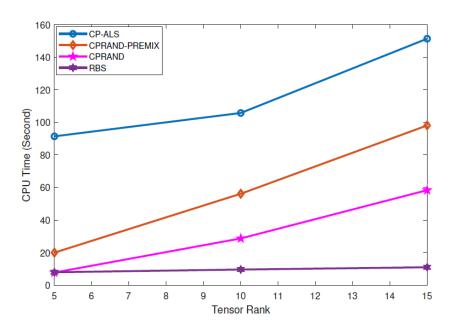
Picture from: E.E. Papalexakis, C. Faloutsos, N.D. Sidiropoulos, *Tensors for data mining and data fusion: Models, applications, and scalable algorithms*, ACM Trans. on Intelligent Systems and Technology, vol. 8, no. 2, 16:1–16:44, 2016.

$$\mathbf{A}^{(1)} = \operatorname{randn}(500, R), \quad \mathbf{A}^{(2)} = \operatorname{randn}(500, R),$$

$$\mathbf{A}^{(3)} = \operatorname{randn}(500, R). \qquad \eta = 0.01$$

$$\underline{\mathbf{X}} = \left[\left[\mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \mathbf{A}^{(3)} \right] \right] \qquad \underline{\hat{\mathbf{X}}} = \underline{\mathbf{X}} + \eta \left(\frac{\|\underline{\mathbf{X}}\|}{\|\underline{\mathbf{Y}}\|} \right) \underline{\mathbf{Y}},$$

$$f = 1 - \frac{\left\| \widehat{\mathbf{X}} - \widetilde{\mathbf{X}} \right\|}{\left\| \widehat{\mathbf{X}} \right\|}$$



Algorithm	R=5	R = 10	R = 15
CP-ALS	0.9900	0.9900	0.9900
CPRAND-PREMIX	0.9897	0.9898	0.9898
CPRAND	0.9893	0.9895	0.9896
RBS	0.9881	0.9897	0.9898
	'	·	

Coil-100 dataset

S. Nene, S. Nayar, and H. Murase, Columbia Object Image Library (COIL-100), Tech. Report CUCS-006-96, Columbia University, 1996.















100 different object classes

72 different angles

The size of the data: $128 \times 128 \times 3 \times 7200$

For the CP-rank R=20 the following results were achieved

	Cpu Time (Second)	
CP-ALS	220	0.686
CPRAND-PREMIX	77.46	0.684

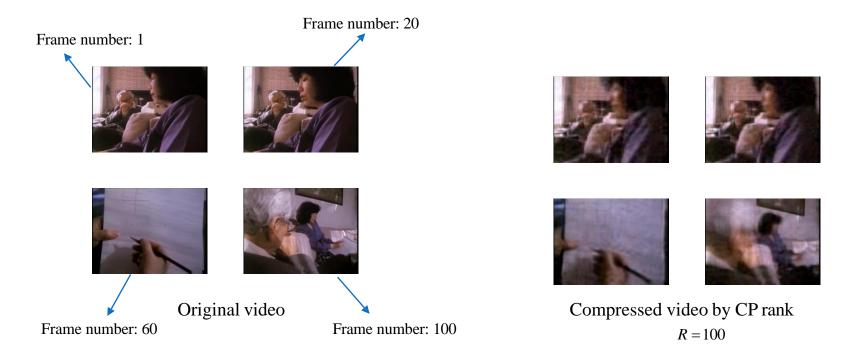
Video compression



Video link: https://www.youtube.com/watch?v=1qeWugmiGt4

The size of this tensors : $360 \times 480 \times 3 \times 2200$

We consider only one and 29 second of that video which is a tensor of size $360\times480\times3\times300$



Here again the CP-PREMIX was 2 times faster than the CP-ALS.

Thanks for your attention!